

## Passing parameters on the stack.

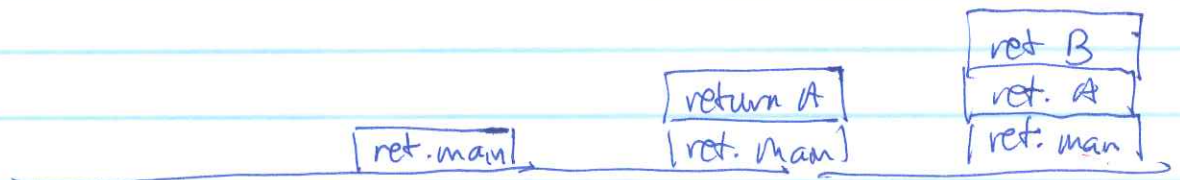
- Short coming of passing parameters in registers:

- one-level ~~sub~~ subroutine call only

- More than one level subroutine call requires subroutine to save parameters to memory locations (very clumsy).

- A better way to pass parameters:

Consider function activation sequence:



↑  
A's parameter is only needed while A is active

↑ ↗  
same holds for B and C

(not needed when A exits - returns to main.)

- Conclusion: using the stack to pass parameters will be very effective, because the parameters are only on the stack as long as the function is active!

- How to pass parameters on the stack.

Example:

```
main: int a, b, c;
      c = sum(a, b)
```

```
int sum(int x, int y)
{
  return (x2 + y2);
}
```

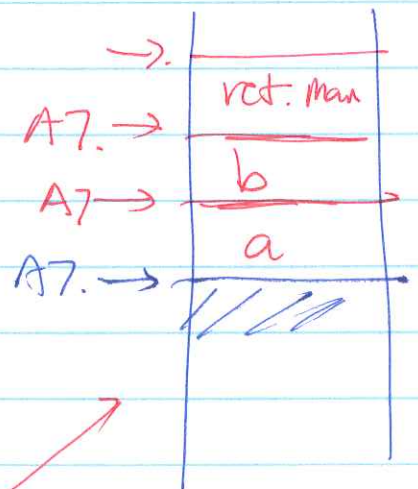
Before:

```
main: l a, d, x
      movl b, d1
      bsr sum.
```

you must pass parameter first

then call subroutine.

↓ pass a, b on stack.



Stack looks like this.

```
subr: l #4, A7
      movl a, (A7)
      subr #4, A7
      movl b, (A7)
```

Special push addressing mode

~~Because~~

The operation to push a value on the system stack is:

- ① suba.l #4, A7 - increase stack top
- ② move.l value, (A7) - save on stack.

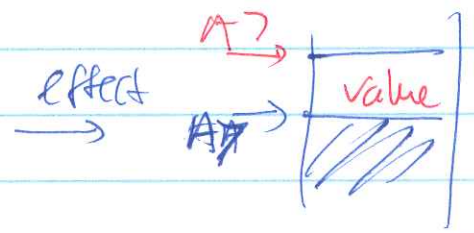
This happens so often, M68000 has a special addressing mode for this:

move.l value, -(A7)

- ① decrement A7 by an amount equal to long (4) (suba.l #4, A7)
- ② move.l value, (A7).

You can also push words:

move.w value, -(A7)



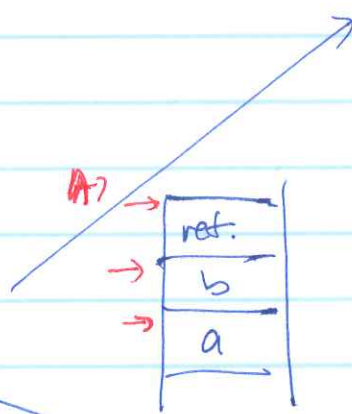
- = ① suba.l #2, A7
- ② move.w value, (A7)



Caller:

move.l a, -(A7)  
move.l b, -(A7)

call sum



Sum:

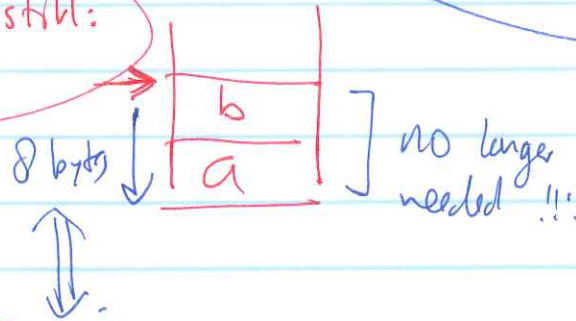
move.l 8(A7), d0  
muls d0, d0

move.l 4(A7), d1  
muls d1, d1

add.l d0, d1  
RTS

only pops  
return address!

Stack is still:



adda.l #8, A7

move.l d0, c

clean w stack

Demo:

sub-stack-param l.s

## Passing Parameters on Stack: Local variables on stack

- For the same reason why we want to pass parameters on the stack, we should put local variables on the stack also:
  - the local variables are only "active" (needed) when the function is running.
  - They must be clean up when the function returns.
  - It is very effective to use the stack to store local variables.

### • Example program:

main:

```
int A[10]
```

```
int sum
```

```
sum = ArraySum(A, 10);
```

```
int ArraySum (int A[],  
              int n)
```

```
{  
  int i, s;
```

```
  s = 0;
```

```
  for (i = 0; i < n; i++)
```

```
    s = s + A[i];
```

```
  return (s);  
}
```

Mam:

Array Sum:

```

move.l #A, -(A7)
move.l #10, -(A7)

```

```

suba.l #8, A7

```

bsr Array Sum

```

adda.l #8, A7
move.l d0, sum

```

stack now:

|        |             |      |
|--------|-------------|------|
| 0(A7)  | (S)         | ← A7 |
| 4(A7)  | (i)         |      |
| 8(A7)  | return addr |      |
| 12(A7) | #10 (n)     |      |
| 16(A7) | #A (A)      |      |

```

move.l #0, (A7) [s=0]

```

```

move.l #0, 4(A7) [i=0]

```

```

while: move.l 4(A7), D0 [i]
       move.l 12(A7), D1 [n]
       cmp.l D1, D0 i ≥ n

```

BGE whileEnd

```

[
  move.l 16(A7), A0
  move.l 4(A7), D0
  muls #4, D0
  move.l (A0, D0.w), D0
  add.l D0, (A7)
]

```

DEMO  
 sub-stack-param2.S

```

whileEnd: move.l (A7), D0

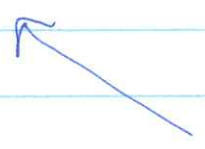
```

```

adda.l #8, A7

```

RTS



```

addq.l #1, 4(A7)

```

BRA WHILE



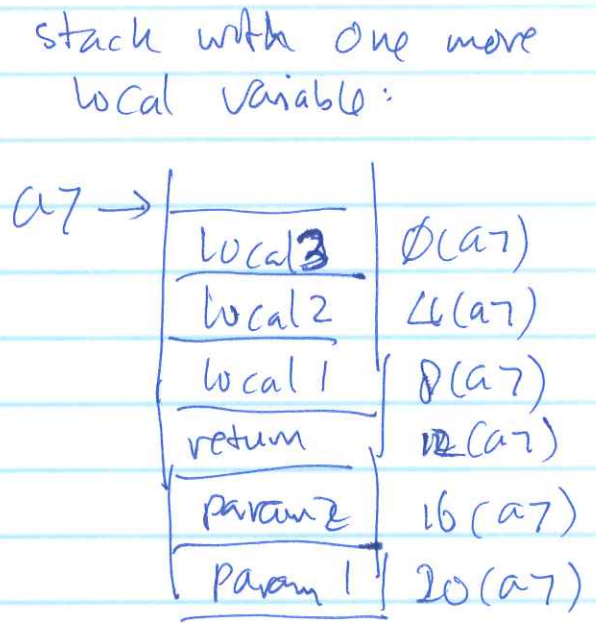
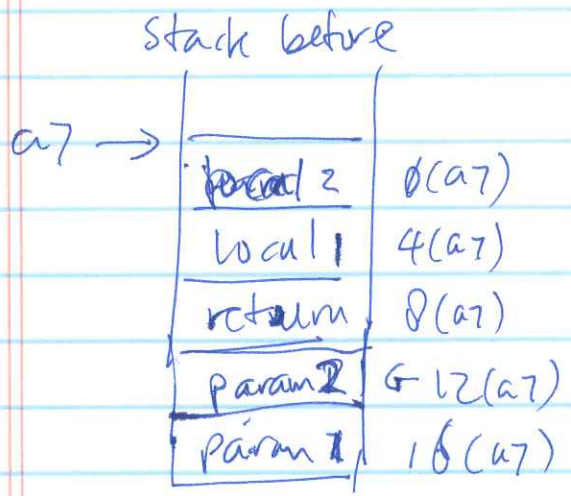
• This solution works, but it is very clumsy:

if at a future time, the algorithm/program need to be expanded with:

- addition parameters
- or - addition local variables (to write program)

then the WHOLE program changes:

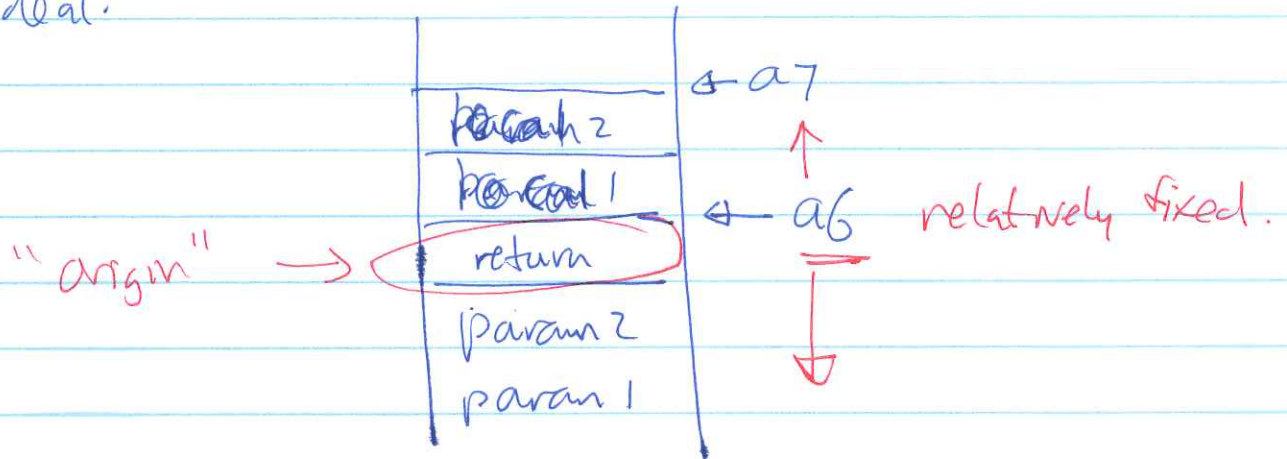
Eg: adding one more local variable:



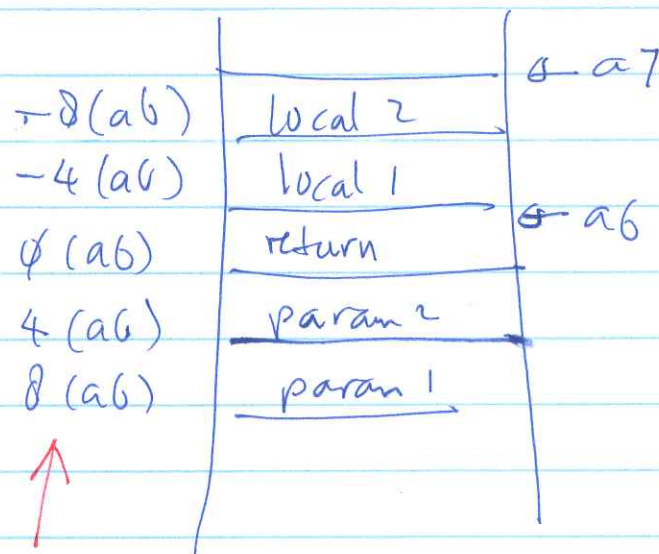
All offsets in the program must be recalculated !!!

Improved solution : use fixed "frame pointer" to access parameters and local variables.

Ideal:



Scheme for accessing local variables and parameters:



neg. offsets to access local variables

positive offsets to access parameters.



main: (unchanged)

```
move.l #A, -(a7)
move.l #0, -(a7)
```

bsr ArraySum

```
adda.l #0, a7
move.l d0, sum
```

ArraySum:

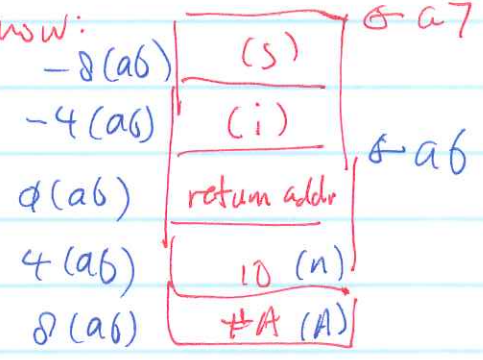
```
movea.l a7, a6
```

make a6 point to base

```
suba.l #0, a7
```

(reserve 2 local vars.)

Stack is now:



Remo  
sub-stack-param3.s

```
move.l #0, -8(a6) [s=0]
move.l #0, -4(a6) [i=0]
```

```
while: move.l -4(a6), D0 [i]
       move.l 4(a6), D1 [n]
       cmp.l D1, D0
```

BGE while\_end

```
movea.l 8(a6), #0
move.l -4(a6), d0
muls #4, d0
move.l (a0, d0, 0), d0
add.l d0, -8(a6)
addq.l #1, -4(a6)
BRA while
```

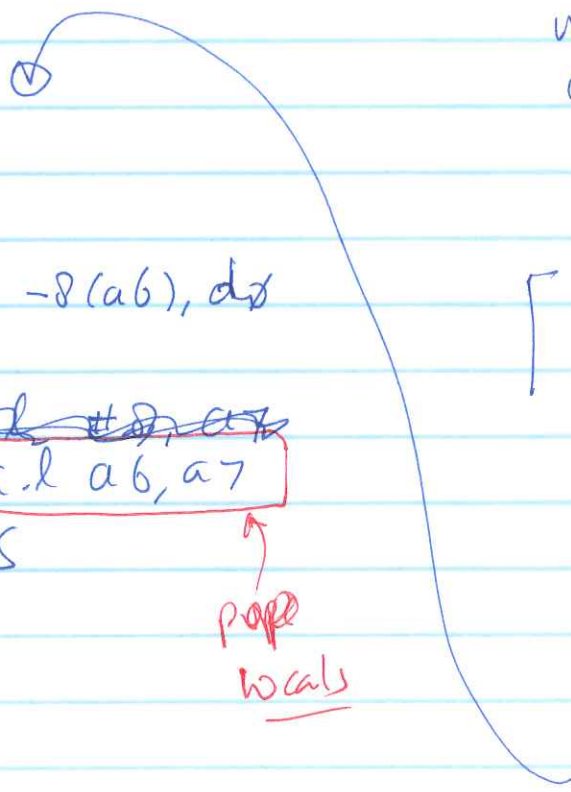
while\_end:

```
move.l -8(a6), d0
```

```
adda.l #0, a7
movea.l a6, a7
```

RTS

pop  
locals

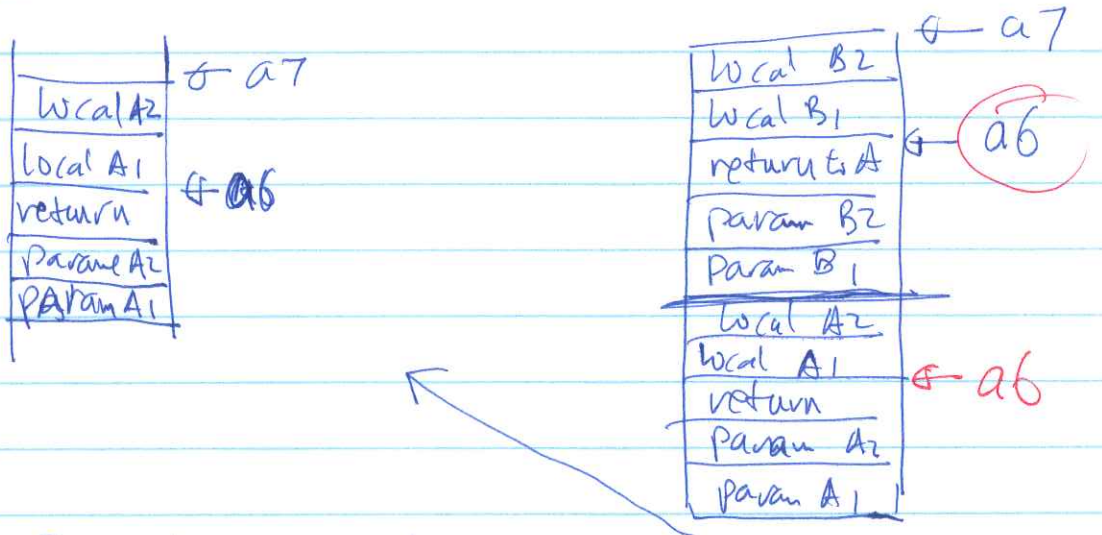


There is ONE problem with the solution:

It does not work with multiple level subroutine call.

Reason:

main → A → B.



when B returns to A.

(if successful)

we lost the value of a6!!!

There is no way to recover the value of a6!

Solution: save the value of a6

restore value of a6 when we return!!!

Question: Where ~~how~~ should we save a5 of caller?

Answer: Stack!!!

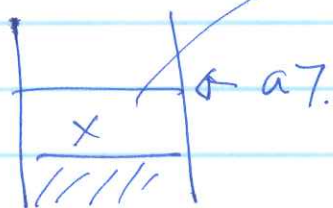
Because the saved value is only needed as long as the subroutine is active.

When subroutine returns, we restore the caller's a5 value, and the saved value is no longer needed!

This requires popping from stack.

How to pop off a stack:

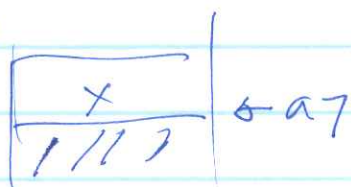
Before:



① `move.l (a7), dp`

② `adda.l #4, a7`

After:



New instruction mode:

`move.l (a7)+, dp`



(1) move.l (a7)+, dn

$\equiv \begin{cases} \text{move.l} & (a7), dn \\ \text{adda.l} & \underline{\#4}, a7 \end{cases}$

(2) move.w (a7)+, dn

$\equiv \begin{cases} \text{move.w} & (a7), dn \\ \text{adda.l} & \underline{\#2}, a7 \end{cases}$

main: (unchanged)

move.l #A, -(a7)  
move.l #10, -(a7)

bsr ArraySum

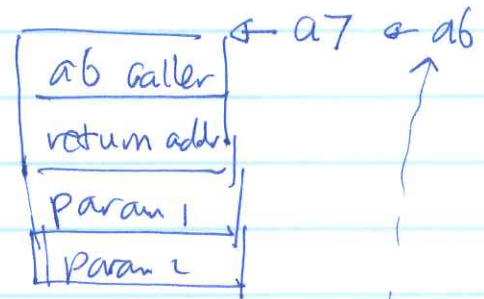
adda.l #0, a7  
move.l d0, sum

ArraySum:

move.l a6, -(a7)

Save a6  
first

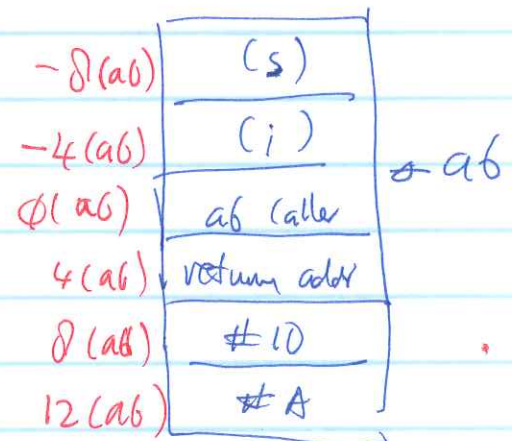
stack now:



movea.l a7, a6

suba.l #0, a7.

stack now:



rest of program same.

make sure restore a6!

Continued



```

movl  #0, -8(a6)    [i=0]
movl  #0, -4(a6)    [i=0]

```

while:

```

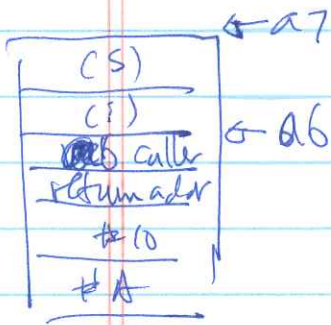
movl  -4(a6), d0    [i]
movl  -8(a6), d1    [n]

cmpl  d1, d0
BGE  while_end

```

Demo

sub-stack-params



```

[
  movl  12(a6), a0
  movl  -4(a6), d0
  mull  #4, d0
  movl  (a0, d0.w), d0
  addl  d0, -8(a6)
]

```

addq.l #1, -4(a6)

BRA while

while\_end:

movl -8(a6), d0 return(s)

movl a6, a7

movl (a7)+, a6

RTS.

Cover a6 for caller !!!

