

## Subroutines with Parameters

- Parameters are used to pass information between the caller and the callee.
- The way that the caller and callee pass input parameters and output (result) to each other is through a rigid agreement (contract).
- Caller ~~agree~~ agrees to put the <sup>input</sup> parameters in a specific location, and the callee agrees to get the input parameters from that location.
- Similarly, callee agrees to put return values in a specific location, and the caller agrees to get the return values from that ~~return~~ location.

Romeo  $\leftrightarrow$  Juliet  
- meeting place agreement.

Q: What can you use to pass parameters?

A: 3 places:

(1) registers. — non-recursive subroutines

(2) memory (rarely used) — non-recursive

(3) stack !!! — recursive subroutines.

Special place in memory.

- We will first look at non-recursive subroutines and only at how parameters are passed in registers.

Let's look at the different ways to pass  
parameters in registers

— by value

↘ by reference.

```
main()
```

```
{ int a, b, c, d;
```

```
  add a, b and put sum in 'sum'
```

```
  add c, d and put sum in 'sum'
```

```
}
```

```
Add(x, y)
```

```
{ add x, y and return result in sum }
```

There are many ways to accomplish this.

Method 1:

subl.s  
by value

```
main()
{ int a, b, c, d, sum;

  sum = Add (a, b);
  sum = Add (c, d);
}
```

```
Add (int x, int y)
{ return (x + y);
}
```

Do Add FIRST

```
f()
{ int c, d, sum;
  sum = Add (c, d);
}
```

a & b are passed by value

i.e.: caller passes the value of the parameters to the callee.

Key: make agreement on:

- (1) where to put param 1, 2, 3, ... etc
- (2) How params are passed.



```

*
* sub1.s: Add subroutine
*     input: d0 = first number
*           d1 = second number
*     output: d0 = sum
*
xdef start, stop, pause1
xdef a, b, c, d, sum

```

*demo  
sub1.s*

```

start:
    move.l  a, d0          ; put first parameter in agreed location
    move.l  b, d1          ; put second parameter in agreed location
    jsr    Add
    move.l  d0, sum        ; get sum from agreed return location

pause1:
    move.l  c, d0          ; put first parameter in agreed location
    move.l  d, d1          ; put second parameter in agreed location
    jsr    Add
    move.l  d0, sum        ; get sum from agreed return location

stop:   nop

a:      dc.l 23
b:      dc.l 12
c:      dc.l 11
d:      dc.l 67
sum:    ds.l 1

```

```

* ----- This is the only way to show clearly
* ----- where a subroutine starts !!!!!!!!
* Subroutine Add
* Input: d0 = first integer
*       d1 = second integer
* Output: d0 = sum
* ----- If that's not clear enough, I don't
* ----- know what else to do.....

```

```

Add:
    add.l  d1, d0          ; Add inputs and put result in agreed location
    rts

end

```

```

.gdbinit:
break start
break stop
break pause1
disp/dw &sum
disp/dw &d
disp/dw &c
disp/dw &b
disp/dw &a
disp/i $pc

```

Method 2:

Sub | a.S  
Sub | b.S

by reference

```
main ()  
{ int a, b, c, d, sum  
  
  sum = Add (&a, &b);  
  sum = Add (&c, &d);  
}
```

```
Add ( int *x, int *y )  
{ return (*x + *y);  
}
```

```
f() { int c, d, sum  
    sum = Add(c, d);  
}
```

Variables are passed by reference

ie: you pass the address of the variables to the callee.

```

*
* sub2a.s: Add subroutine
*       input: d0 = ADDRESS of the first integer
*       d1 = ADDRESS of the second integer
*       output: d0 = sum
*
xdef start, stop, pause1
xdef a, b, c, d, sum

```

*sub 2a.s*  
*add*  
~~*add*~~  
~~*add*~~

```

start:
    move.l #a, d0 ; put first parameter in agreed location
    move.l #b, d1 ; put second parameter in agreed location
    jsr    Add
    move.l d0, sum ; get sum from agreed return location
pause1:
    move.l #c, d0 ; put first parameter in agreed location
    move.l #d, d1 ; put second parameter in agreed location
    jsr    Add
    move.l d0, sum ; get sum from agreed return location
stop:   nop

a:      dc.l 23
b:      dc.l 12
c:      dc.l 11
d:      dc.l 67
sum:    ds.l 1

```

*not a good place....*

```

* ----- This is the only way to show clearly
* ----- where a subroutine starts !!!!!!!!
* Subroutine Add
* Input: d0 = ADDRESS of the first integer
*       d1 = ADDRESS of the second integer
* Output: d0 = sum
* ----- If that's not clear enough, I don't
* ----- know what else to do.....

```

```

Add:
    movea.l d0, a0 ;
    move.l (a0), d7 ; get first number
    movea.l d1, a0 ;
    add.l (a0), d7 ; add second number to the first
    move.l d7, d0 ; put sum in agreed return location
    rts
end

```

*needless.*

```

.gdbinit:
break start
break stop
break pause1
disp/dw &sum
disp/dw &d
disp/dw &c
disp/dw &b
disp/dw &a
disp/i $pc

```

```

*
* sub2b.s: Add subroutine
*      input: A0 = ADDRESS of the first integer
*            A1 = ADDRESS of the second integer
*      output: d0 = sum
*

```

*demo/sub2b.s*

```

xdef start, stop, pause1
xdef a, b, c, d, sum

```

```

start:
  move.l #a, a0 ; put first parameter in agreed location
  move.l #b, a1 ; put second parameter in agreed location
  jsr    Add
  move.l d0, sum ; get sum from agreed return location

```

```

pause1:
  move.l #c, a0 ; put first parameter in agreed location
  move.l #d, a1 ; put second parameter in agreed location
  jsr    Add
  move.l d0, sum ; get sum from agreed return location

```

```

stop:  nop

```

```

a:    dc.l 23
b:    dc.l 12
c:    dc.l 11
d:    dc.l 67
sum:  ds.l 1

```

```

* ----- This is the only way to show clearly
* ----- where a subroutine starts !!!!!!!!
* Subroutine Add
* Input: a0 = ADDRESS of the first integer
*        a1 = ADDRESS of the second integer
* Output: d0 = sum
* ----- If that's not clear enough, I don't
* ----- know what else to do.....

```

```

Add:
  move.l (a0), d7 ; get first number
  add.l  (a1), d7 ; add second number to the first
  move.l d7, d0  ; put sum in agreed return location
  rts
end

```

*easier.*

```

.gdbinit:
break start
break stop
break pause1
disp/dw &sum
disp/dw &d
disp/dw &c
disp/dw &b
disp/dw &a
disp/i $pc

```



Method 3:

```
main ( )  
{ int a, b, c, d, sum;
```

~~sum~~

```
Add (a, b, &sum);
```

```
Add (c, d, &sum);
```

```
}
```

```
Add ( int x, int y, int *sum)
```

```
{ *sum = a + b; }
```

a & b are passed by value  
sum is passed by reference.

```

*
* sub3.s: Add subroutine
*   input: d0 = the first integer
*          d1 = the second integer
*          a0 = address of the sum variable
*   output: None. Add will update the sum variable
*

```

*ctlimp/sub3.s*

```

xdef start, stop, pause1
xdef a, b, c, d, sum

```

*agreed, not by good place.*

```

start:
    move.l  a, d0      ; put first parameter in agreed location
    move.l  b, a1 d1 ; put second parameter in agreed location
    move.l  #sum, a0   ; put third parameter in agreed location
    jsr     Add

```

```

pause1:
    move.l  c, d0      ; put first parameter in agreed location
    move.l  d, d1      ; put second parameter in agreed location
    move.l  #sum, a0   ; put third parameter in agreed location
    jsr     Add

```

```

stop:  nop

```

```

a:     dc.l 23
b:     dc.l 12
c:     dc.l 11
d:     dc.l 67
sum:   ds.l 1

```

```

* ----- This is the only way to show clearly
* ----- where a subroutine starts !!!!!!!!
* Subroutine Add
* Input: d0 = first integer
*        d1 = second integer
*        a0 = address of sum variable
* ----- If that's not clear enough, I don't
* ----- know what else to do.....

```

```

Add:
    add.l  d1, d0      ; add second number to the first
    move.l d0, (a0)    ; update the sum
    rts

```

end

```

.gdbinit:
break start
break stop
break pause1
disp/dw &sum
disp/dw &d
disp/dw &c
disp/dw &b
disp/dw &a
disp/i $pc

```

## Method 4

```
main ( )
```

```
{ int a, b, c, d, sum;
```

```
    Add ( &a, &b, &sum );
```

```
    Add ( &c, &d, &sum );
```

```
}
```

```
Add ( int *x, int *y, int *sum )
```

```
{ *sum = *x + *y }
```

```

*
* sub3.s: Add subroutine
*      input: A0 = ADDRESS of the first integer
*             A1 = ADDRESS of the second integer
*             A2 = ADDRESS of the sum variable
*      output: None. Add will update the sum variable
*
      xdef start, stop, pause1
      xdef a, b, c, d, sum

start:
      move.l  #a, a0          ; put first parameter in agreed location
      move.l  #b, a1          ; put second parameter in agreed location
      move.l  #sum, a2        ; put third parameter in agreed location
      jsr     Add

pause1:
      move.l  #c, a0          ; put first parameter in agreed location
      move.l  #d, a1          ; put second parameter in agreed location
      move.l  #sum, a2        ; put third parameter in agreed location
      jsr     Add

stop:   nop

a:      dc.l 23
b:      dc.l 12
c:      dc.l 11
d:      dc.l 67
sum:    ds.l 1

```

```

* ----- This is the only way to show clearly
* ----- where a subroutine starts !!!!!!!!
* Subroutine Add
* Input: a0 = ADDRESS of the first integer
*        a1 = ADDRESS of the second integer
*        a2 = ADDRESS of the sum variable
* ----- If that's not clear enough, I don't
* ----- know what else to do.....

```

```

Add:
      move.l  (a0), d7        ; get first number
      add.l   (a1), d7        ; add second number to the first
      move.l  d7, (a2)        ; update the sum
      rts

      end

```

```

.gdbinit:
break start
break stop
break pause1
disp/dw &sum
disp/dw &d
disp/dw &c
disp/dw &b
disp/dw &a
disp/i $pc

```



## Subroutines with local variables (Non-recursive).

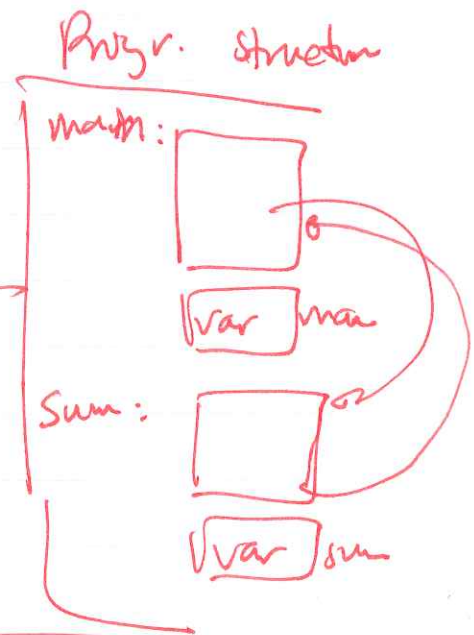
Example: write a subroutine that ~~sum~~ returns the sum of an array A of N elements.

```
int sum (int A[], N)
{
    int result, i;
    result = 0;
    for (i = 0; i < N; i++)
        result = result + A[i];
    return (result);
}
```

```
main ()
{
    int B[10];
    int S;

    S = sum (B, 10);
}
```

we branch  
Programs  
to run  
in legal  
places!



demo/local-variables

Decision:

- (1) **A** is passed by reference. in d0
- (2) **N** is passed by value in d1
- (3) result is returned in d0.

- The way that main calls sum is then as follows:

```

move.l #B, d0
move.l #10, d1
bsr    sum
    } sum(B, 10)

```

```

move.l d0, s
    } put return value
      in s.

```

- The subroutine sum has 2 local variables:

result and i.

We can define these variables as:

```
result: ds.l 1
```

```
i: ds.l 1
```

- program using these variables.

Alternately, we can associate with each variable a register and program the subroutine sum to use registers

- Warning: reserving storage for local variables using ds or registers only works with NON-RECURSIVE functions.

```

* sub5.s
*
* Illustrates local variables in non-recursive functions.
* Here, we put local variables in memory.
*
        xdef start, stop, B, s, result, i
        suba.l #80, SP          ; So I don;t crash when displaying the stack
* ----- Main
* main()
* { int B[10], s;
*   s = sum(B, 3);
* }
* -----
start:
        move.l #B, d0
        move.l #10, d1
        bsr    sum              ; Call sum with B and 3
        move.l d0, s            ; Put return value in s
stop:   nop

B:      dc.l 2, 1, 3, 4, 1, 2, 1, 1, 2, 2
s:      ds.l 1

* ***** Subroutine sum with local variables in memory
* int sum(int A[], int N)    ==> A[] in d0, N in d1
*                          ==> result in d0
* { int result
*   int i
*
*   result = 0;
*   for (i = 0; i < N; i++)
*     result += A[i];
*   return result;
* }
* -----
* Note: d0 = address of array A[]
*       d1 = value of N
* *****
sum:    move.l #0, result
        move.l #0, i

forloop:cmp.l i, d1          ; Test if i >= N
        ble   fordone       ; If so, exit for-loop

        move.l d0, a0        ; Get start of array in a0
        move.l i, d2         ; Compute offset
        muls  #4, d2         ;
        move.l 0(a0,d2.w), d2 ; Fetch A[i] in d2
        add.l d2, result     ; result := result + A[i]

        addq.l #1, i        ; increment i
        bra   forloop

fordone:move.l result, d0    ; Return subprogram return value in d0
        rts

* ----- Local variables for sum subprogram
result: ds.l 1
i:      ds.l 1
* ----- End of sum subroutine
        end

```

The sum subroutine with local variables in memory.

sum:      move.l    #0, result      | d0 = A  
   | d1 = N

            move.l    #0, i

forloop:    move.l    i, d7  
            cmp.l    d7, d7      (d1 = N) (done if:  $i \geq N$ )  
            bge     fordone

~~move.l    i, d7~~  
            move.l    d0, a0  
            move.l    i, d7  
            mul.l    #4, d7  
            move.l    0(a0, d7), d7

            move.l    result, d6  
            add.l    d6, d7  
            move.l    d7, result

            move.l    i, d7  
            add.l    #1, d7  
            move.l    d7, i  
            bra      forloop

fordone:    move.l    result, d0      (return result in d0)  
            rts

result:     ds.l     1  
i:          ds.l     1



## The sum subroutine with local variables in registers.

We use d7 for result  
and d6 for i.

```
sum:  move.l  #0, d7
```

```
      move.l  #0, d6
```

```
forloop:
```

```
move.l
```

```
      cmp.l   d1, d6
```

```
      bge    fordone
```

```
      move.l  d0, a0
```

```
      move.l  d6, d5
```

```
      mult   #4, d5
```

```
      move.l  0(a0, d5), d4
```

```
      add.l   d4, d7
```

```
      add.l   #1, d6
```

(i = i + 1)

```
      bra    forloop
```

```
fordone:  move.l  d7, d0
```

(return result in d0)

```
      rts .
```

Warning: Before you branch to a subroutine, make sure you save all important temporal results in registers to their corresponding variables.

eg: Because if you don't, ~~the~~ subroutine and the subroutine uses some of the registers for calculation, then your temporal results will be overwritten and ~~not~~ no longer available.

• Subroutines with ~~parameters~~ local variables in stack:

