

4

Real Numbers

Real Number



Can't be represented precisely

Only approximations



Fixed point numbers

Floating Point numbers

- Before we discuss the representation of fixed and floating point numbers, we look at how fractions are converted between decimal & binary number systems.

4

Converting fractions

- Binary \rightarrow Decimal

$$(0.01101)_2 = 0.01101$$

$0 \quad -1 \quad -2 \quad -3 \quad -4 \quad -5$
 $2^{-2} \quad 2^{-3} \quad 2^{-5}$

$$\frac{1}{4} + \frac{1}{8} + \frac{1}{32} = \frac{8+4+1}{32} = \frac{13}{32}$$

$$= (0.40625)_{10}$$

- Decimal \rightarrow Binary : multiply by 2 & write integer part.

$$(0.6875)_{10} = (\dots)_2$$

$\begin{array}{r} \times 2 \\ \hline 0.6875 \\ \hline 1.3750 \\ \text{take remainder } \underline{\underline{0}} \end{array}$

 $\begin{array}{r} \times 2 \\ \hline 0.3750 \\ \hline 0.750 \\ \text{take remainder } \underline{\underline{0}} \end{array}$

 $\begin{array}{r} \times 2 \\ \hline 0.750 \\ \hline 1.5 \\ \text{take remainder } \underline{\underline{1}} \end{array}$

 $\begin{array}{r} \times 2 \\ \hline 0.5 \\ \hline 1.0 \\ \text{take rem. } \underline{\underline{0}} \end{array}$

done.

answer.

$$(0.1011)_2$$

ans: $0.011001100110\dots$

repeat  repeats.

read.

$$(0.1)_{10} = (\dots)_2$$

$$\begin{array}{r}
 \times 2 \quad 0.4 \\
 \hline
 0.8 \\
 \times 2 \quad 1 \\
 \hline
 1.6 \\
 \hline
 0.6 \\
 \times 2 \quad 1 \\
 \hline
 1.2 \\
 \hline
 0.2 \\
 \times 2 \quad 0 \\
 \hline
 0.4
 \end{array}$$

- Decimal \rightarrow octal multiply by 8, note integer part.

$$(0.12)_{10} = (\dots)_8 = (0.0753\dots)_8$$

$$\begin{array}{r} 0.12 \\ \times 8 \\ \hline 0.96 \end{array} \quad 0$$

$$\begin{array}{r} 0.96 \\ \times 8 \\ \hline 7.68 \end{array} \quad 7.$$

$$\begin{array}{r} 0.68 \\ \times 8 \\ \hline 5.44 \end{array} \quad 5.$$

$$\begin{array}{r} 0.44 \\ \times 8 \\ \hline 3.52 \end{array} \quad 3$$

etc.

Fixed point ~~not~~ numbers

- integer : $(10101010)_2 = 2^7 + 2^5 + 2^3 + 2^1 = 128 + 32 + 8 + 2 = \underline{\underline{170}}$

It is assumed that the decimal point lies at the right most digit :

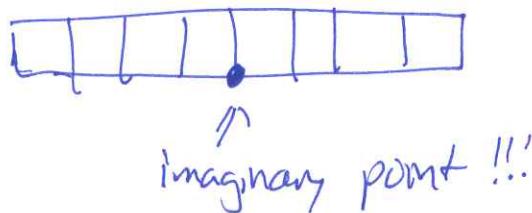
10101010.

- Fixed point number:

is very similar to integer representation, except the decimal point is assumed to lie elsewhere.

Note: computer CANNOT place a decimal point explicitly.
It must assume the ~~local~~ ~~a~~ fix place for it.

e.g.: Assume decimal point lies at :



Then the ~~number~~ bit pattern

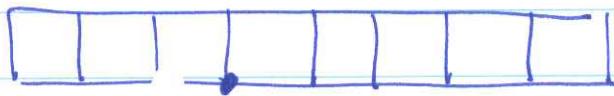
10101010

means: $\begin{array}{ccccccc} 3 & 2 & 1 & 0 & -1 & -2 & -3 & -4 \\ 1010.1010 \end{array}$

$$8 + 2 + \frac{1}{2} + \frac{1}{8} = 10 \cdot \underline{\underline{\frac{5}{8}}} = 10.625$$

Converting Fixed Point Binary numbers.

Assume binary fixed point representation is:



Binary to decimal conversion:

Pattern: $\boxed{01101100}$

This implies (assume decimal point):

011.01100

$2^1 2^0 2^{-1} 2^{-2} 2^{-3}$

$$\text{Value} = 2 + 1 + \frac{1}{4} + \frac{1}{8} \cdot$$

$$= 3 \frac{3}{8}$$

$$= 3.375$$

Decimal to binary (fixed point) conversion:

$$(3.14)_{10} = ?$$

Split 3.14 into integer part $\rightarrow 3$
and fraction part $\rightarrow 0.14$.

Integer part converted as follows:

$$\begin{array}{r} 3 \\ 2 \overline{)1} \\ \underline{-1} \\ 2 \overline{)1} \\ \underline{-0} \end{array} \quad (3)_{10} = (11)_2.$$

Fraction part converted as follows:

$$\begin{array}{r} 0.14 \\ \times 2 \overline{\longrightarrow} 0 \\ 0.28 \end{array} \quad (0.14)_{10} = (0.001000\dots)_2$$

$$\begin{array}{r} 0.28 \\ \times 2 \overline{\longrightarrow} 0 \\ 0.54 \end{array}$$

$$\begin{array}{r} 0.54 \\ \times 2 \overline{\longrightarrow} 1 \\ 1.08 \end{array}$$

$$\begin{array}{r} 0.12 \\ \times 2 \overline{\longrightarrow} 0 \\ 0.24 \end{array}$$

$$\begin{array}{r} 0.24 \\ \times 2 \overline{\longrightarrow} 0 \\ 0.48 \end{array}$$

$$\begin{array}{r} 0.48 \\ \times 2 \overline{\longrightarrow} 0 \\ 0.96 \end{array}$$

so:

$$(3.14)_w = (11.001000\dots)$$

From 8 bits and line up the decimal point:

.	1	1	0	0	1	0	0
---	---	---	---	---	---	---	---

↑

answer.

Conclusion: ~~Real numbers are NOT represented precisely.~~

- Integer numbers can be represented precisely by computers.
- Non-integer numbers CANNOT.
There will be a round off / truncation error for non-integer numbers.
The magnitude of the error depends on the number of bits used to represent the fraction.

Representing Floating Point numbers

- Floating point \neq Real !

Real number can have infinite number of digits.

- Floating point number n :

$$n = m \times 10^e$$

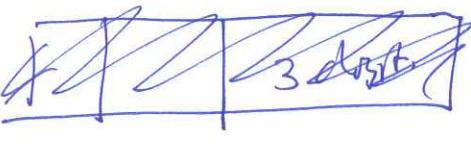
mantissa exponent.

(for humans).

$$\begin{aligned} \text{eg: } 19.47 &= 0.1947 \times 10^2 \\ &= 1947.0 \times 10^{-2}. \quad \text{etc.} \end{aligned}$$

- ~~Not every possible number~~

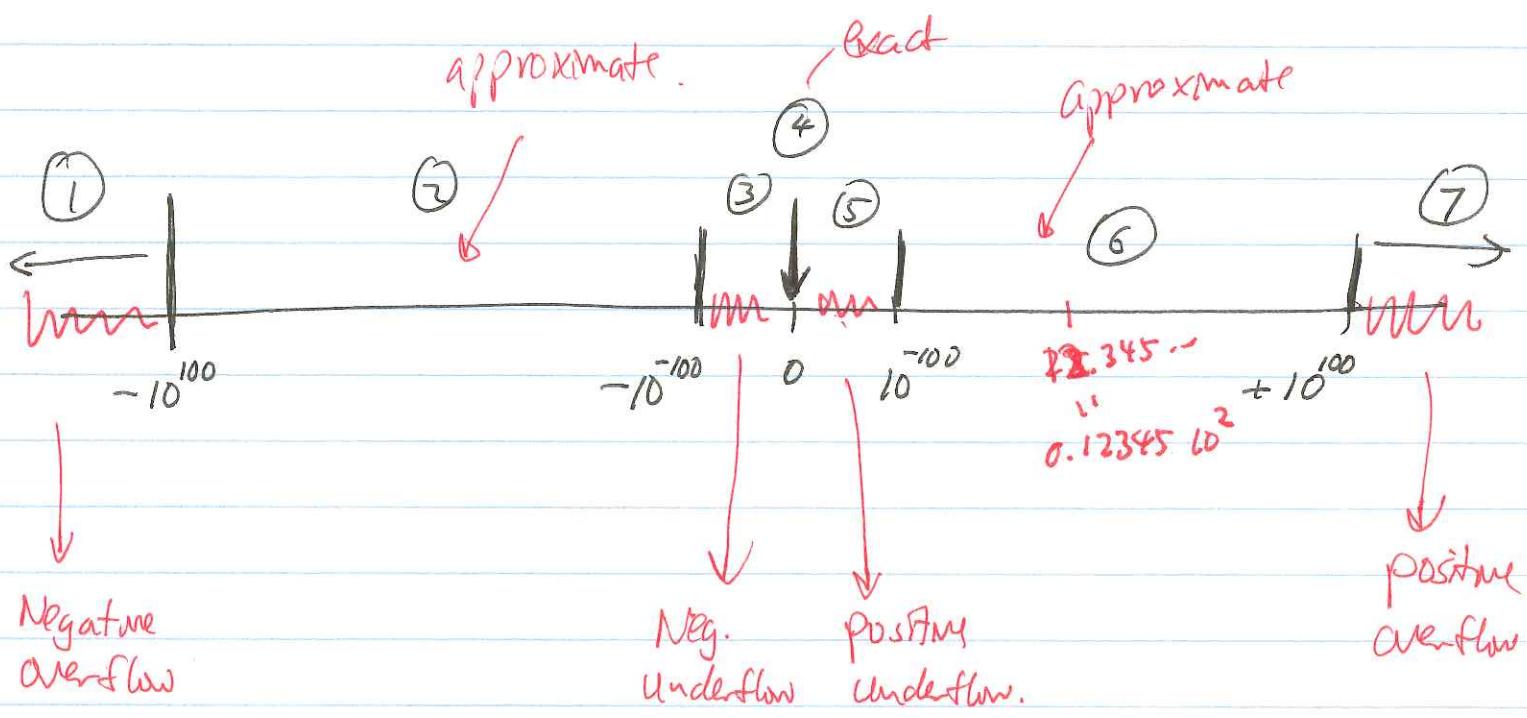
- Mantissa & exponent are of finite length.
- Not every possible numerical value can be expressed.

Eg. Si 

decimal

Example: 3 digit mantissa of |exponent| ≤ 99

The real numbers ~~then~~ can be divided into 7 regions:



(~~Too big~~
Exponent
too big).

(exponent
too small)

- 10

Any number between $(-1 \cdot 10^{100}, -1 \cdot 10^{-10})$.

IEEE floating-point standard 754

- IEEE 754 defines 3 standard formats for floating point numbers.

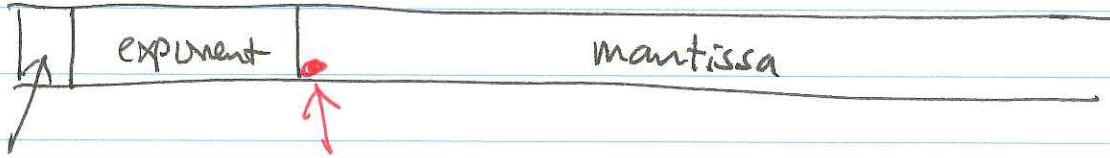
(Most computers now use these formats).

- single precision. (32 bits)
- double precision (64 bits)
- Extended precision (80 bits)

- Double precision format: (others are similar)

bits 1 11

52



sign.
mantissa.

decimal point is here.

- the mantissa is between 1 & 2.:

$1 \leq m < 2$. (Decimal point assumed HERE)

- Exponent is (2^e) in excess 2^{10} .

(Thus: the mantissa is signed magnitude.

the exponent is excess 2^{10}).

Operations with floating point numbers:

- multiply & divide :
 - (1) multiply / divide mantissa
 - (2) add / subtract exponent
 - (3) Normalize !
 - bring result back to a form that $1 \leq \text{mantissa} < 2$
- add & subtract :
 - (1) bring both operands to same exponent value.
 - (2) add / subtract mantissa
 - (exponent = exponent of one operand)
 - (3) Normalize.

Floating Point Number

- Shortcoming of fixed point number : bound by range
- A better way to represent real number is the "scientific" notation :

$$\text{real number} = \text{fraction} \times 10^{\text{exponent}}$$

↑

fraction is also called mantissa

- The computer version of this way of representing real numbers is called "floating point".
- Example :

$$314 = 3.14 \times 10^2 = 0.314 \times 10^3$$
$$0.000012 = 1.2 \times 10^{-5} = 0.12 \times 10^{-4}$$

- There are many ways to represent the same number :

$$314 = 3.14 \times 10^0 = 0.314 \times 10^1 \text{ etc.}$$

^{↑ form 1.}

One form is usually chosen as the standard .
(ie: where to put the decimal point).

- Standard forms typically chosen are:

(1) mantissa \neq consists of 1 digit ($\neq \emptyset$)
number stored as:

eg:	$314 = 3.14 \times 10^2$	$31400 +02$
	$31.4 = 3.14 \times 10^1$	$31400 +01$
	$1255 = 1.255 \times 10^3$	$12550 +03$
	$0.00012 = 1.2 \times 10^{-4}$	$12000 -04$

(2) mantissa is \emptyset and leading digit of fraction must not be \emptyset .

number stored as

eg:	$314 = 0.314 \times 10^3$	$31400 +03$
	$31.4 = 0.314 \times 10^2$	$31400 +02$
	$1255 = 0.1255 \times 10^4$	$12550 +04$
	$0.00012 = 0.12 \times 10^{-3}$	$12000 -03$

- Now do this in binary:

~~$1011 \times 2 = 1100$~~

$$= 10.1 \times 2 = 1100_5$$

Note that:

101	$= 10.1 \times 2^1$
1	$= 1.01 \times 2^2$
"5"	$= 0.101 \times 2^3$

There also, there are many way to represent the same floating point number.

One way is picked as the standard.

- Eg: mantissa consists of 1 digit ($\neq 0$)

Then:

<u>number</u>	<u>standard repr</u>	<u>internally stored as</u>
101	$= 1.01 \times 2^2$	10100000 0010
0.101	$= 1.01 \times 2^{-1}$	10100000 "0001"
-101	$= -1.01 \times 2^2$	"-10100000" 0010
-0.101	$= -1.01 \times 2^{-1}$	"-10100000" "0001"

Sj: mantissa AND exponent can be pos & neg. numbers ???

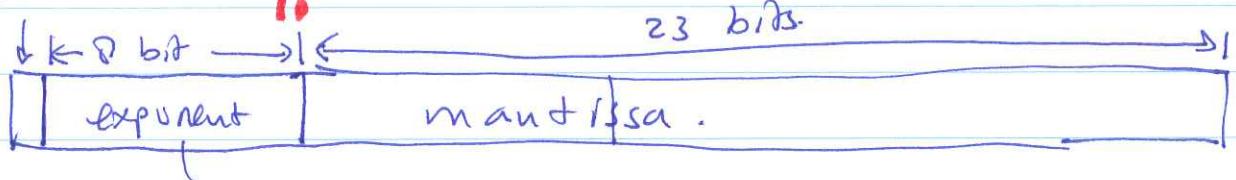
- Single precision floating point number : 4 bytes.

Standard form : 1 digit mantissa ($\neq \phi$).

Example :

$$101 = 1.01000 \dots \times 2^2$$

sign of mantissa \downarrow Assumed 1. . .



use	1.0	1.5	1.75
	2.0	3	3.5
	4.0	6	7

- Mantissa is stored in Sgn-magnitude representation
unsigned value
or absolute value.

Where the assumed decimal point is immediately after the leading digit.

(mantissa is always 1.)

Also: the leading 1 in mantissa is NOT stored, but it is assumed (since it's always there!)

- Exponent is stored in 8 bits with excess 127 encoding (i.e.: value is stored as absolute (value+127))
- Example:

$$(1.0)_0 = (1.0)_2 \times 2^0$$

Mantissa = 1.00000.....0
 23 bits. sign = Ø.

Exponent = 0

representation in excess 127

Value x is repr. by absol. value $x+127$

0 is repr. by ~~as~~: unsigned value

$$0+127 = 127 = 0111.1111 \\ (8 \text{ bits})$$

Put it in the format specified:



Example: $(-1.0)_{10}$

$$(-1.0)_{10} = (-1.0)_2 \times 2^0.$$

Mantissa = sign = 1

Magnitude = $1.\underbrace{000\dots}_{23 \text{ bits.}} 0$

Exponent = 0 \rightarrow 127 = 0111.111

Repr:

[1] [0111.111] [000...]

Example: $(1.25)_{10}$.

$$(1.25)_{10} = (1.01)_2$$

$$= (1.01)_2 \times 2^0.$$

Mantissa : sign = 0.

Magnitude = $1.\underbrace{010000\dots}_{23 \text{ bits.}}$

Exponent = 0.

Repr:

[0] [0111.111] [0100000...0]

Example:

$$(16.25)_{10} = (10000.01)_2$$

$$= (1.000001)_2 \times 2^4$$

~~2⁴~~

Mantissa: sign = \emptyset

magnitude = $1.\underbrace{0000010000000}_{23 \text{ bits}}$

Exponent = 4

in excess 127 repr. by unsigned $\emptyset + 127$
 $= 13\emptyset$

$$13\emptyset = 128 + 3$$

$$= 1000.0000 + 1\emptyset$$

$$= 1000.001\emptyset,$$

Representation is:

$\boxed{0}$ $\boxed{1000.001\emptyset}$ $\boxed{100001000\dots 0}$

Given: repr.:

1 0111.1110 110000...0

What is the ~~number~~ value?

Mantissa: sign = 1 \rightarrow neg.

magnitude = 1.110000

$$= \cancel{1} + \cancel{\frac{1}{2}} + \cancel{\frac{1}{4}} = 0.75.$$

So:

~~-1.75~~

Exponent: 0111.1110 = 126

$$\text{repr. value } x + 127 = 126$$

$$x = 126 - 127$$

$$= -1.$$

So:

$$\text{value in binary} = -1.11 \times 2^{-1}$$

$$= -0.111$$

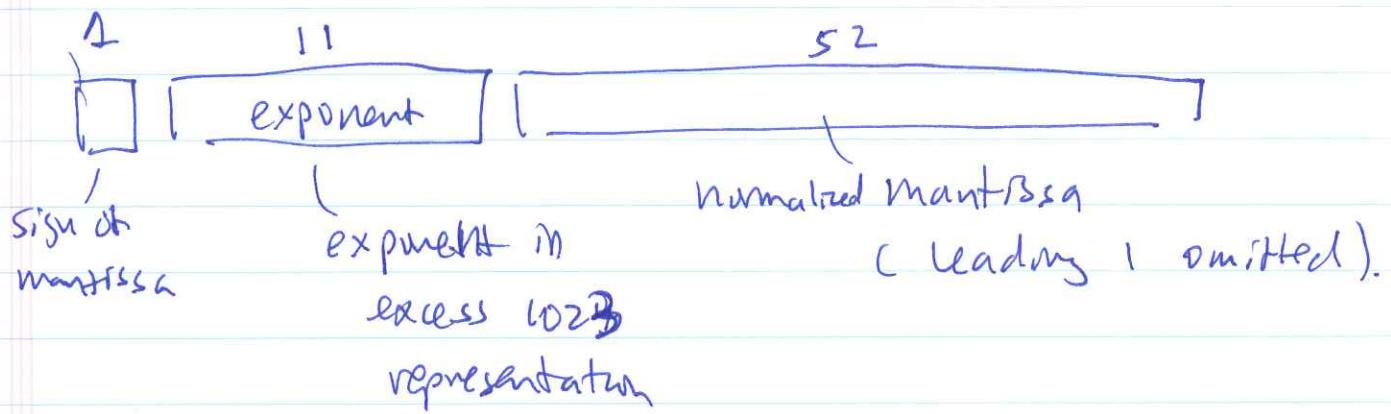
$$= -\left(\frac{1}{2} + \frac{1}{4} + \frac{1}{8}\right)$$

$$= -\frac{7}{8}$$

$$= -0.875$$

=

* Double precision IEEE format :



The IEEE Floating Point Standard

- Up to about 1970, each computer manufacturer has its own way to represent floating point numbers.
- To ease exchange of data (information) between computers, a universal standard for floating point numbers was established by IEEE (Institute of Electrical & Electronics Engineers) in 1985 (The IEEE standard 754)
- Most modern CPUs (Motorola, SPARC, MIPS) include a floating point processor and all of them conform to the IEEE standard.
- The standard defines 3 types of floating point numbers:
 - (1) Single precision (float in C)
 - (2) double precision (double in C)
 - (3) quad precision (long double in C)

They are all similar.