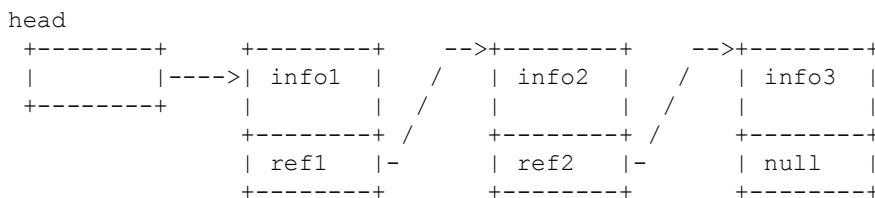
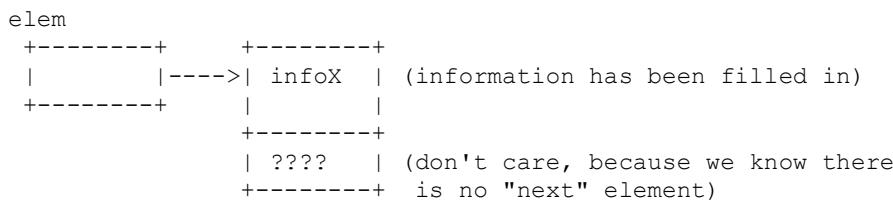


## **Iterative Algorithm to Insert at Tail of a Linked List...**

- Suppose you have a linked list at "head":

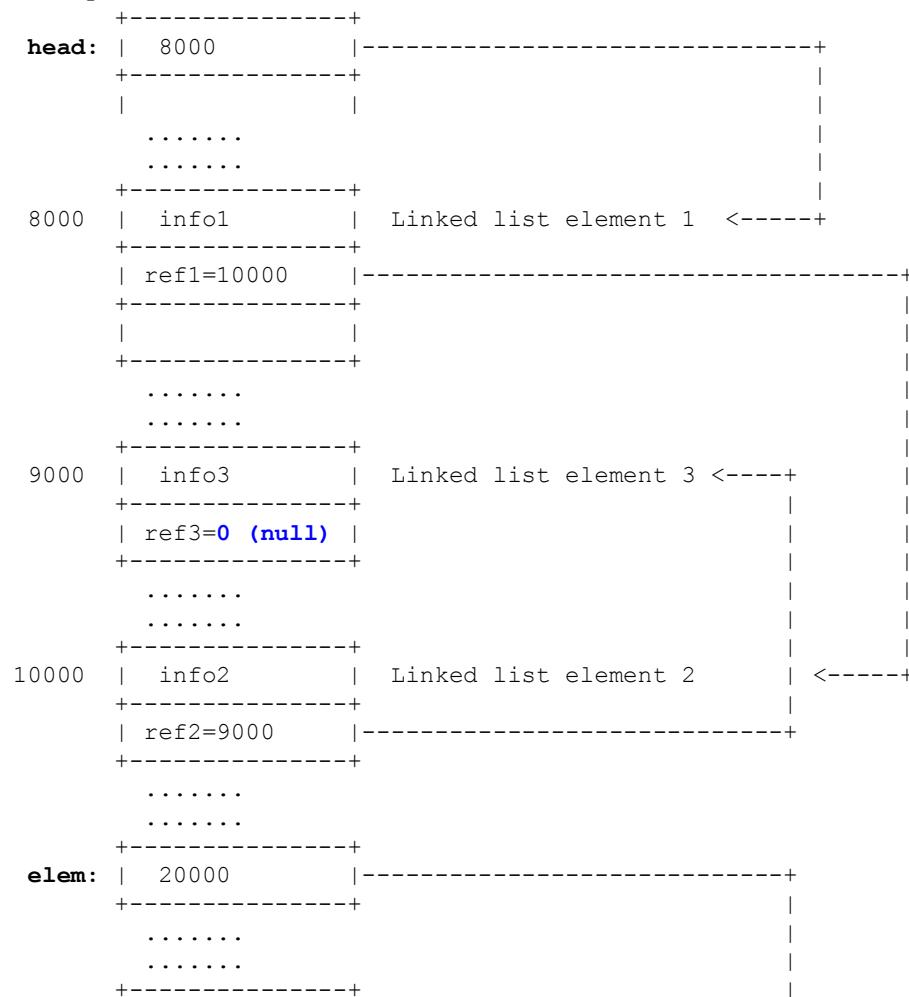


And you have a linked list **element** at "elem":



- Here is an example of how the linked list is stored in memory:

## Memory:



```
20000 | infoX           | New linked list element <----+
+-----+
| refX=????? |
+-----+
```

- Notice that the "links" are memory addresses and the linked list can be found by tracing/following the addresses in the "linkage" field
- Inserting the element **pointed** by "elem" at the **tail** of the linked list is realised by making the last element in the linked list point to the element **pointed** by "elem".

This process is illustrated by the following diagram:

Before insertion:

```
head
+-----+ +-----+ -->+-----+ -->+-----+
|       |---->| info1 | / | info2 | / | info3 |
+-----+ |       | / |       | / |       |
+-----+ / +-----+ / +-----+
| ref1 | - | ref2 | - | null |
+-----+ +-----+ +-----+
elem
+-----+ +-----+
|       |---->| infoX |
+-----+ |       |
+-----+
| ????
+-----+
```

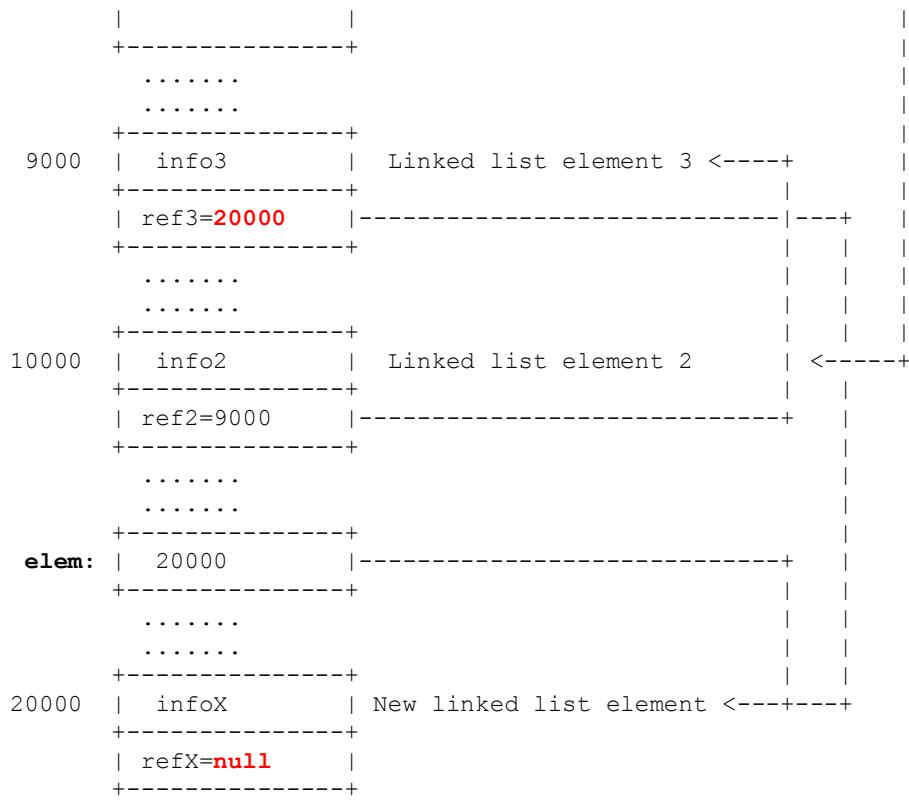
After insertion:

```
head
+-----+ +-----+ -->+-----+ -->+-----+
|       |---->| info1 | / | info2 | / | info3 |
+-----+ |       | / |       | / |       |
+-----+ / +-----+ / +-----+
| ref1 | - | ref2 | - | ref3 |--+
+-----+ +-----+ +-----+ | |
elem
+-----+ +-----+
|       |---->| infoX |<-----+
+-----+ |       |
+-----+
| null |
+-----+
```

- Here is an example of actually happens within the computer memory when you "link" a new element to the list:

Memory:

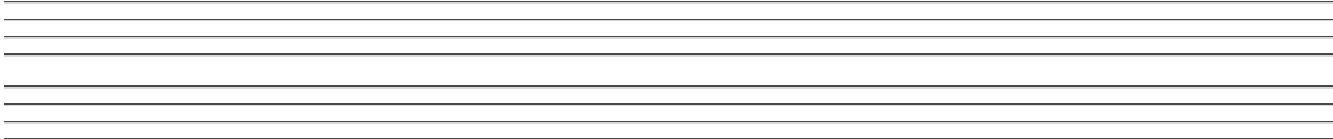
```
+-----+
head: | 8000 |-----+
+-----+ |
|       |
..... |
..... |
+-----+
8000 | infol      | Linked list element 1 <----+
+-----+
| ref1=10000 |
+-----+ |
```



- Notice the change of the address value in the last element of the linked list.

This address (20000) is also the address of the list element pointed to by "elem"

- The other change is the linkage value of the new last element. It must contain ZERO (null) to indicate the end of the list.



- ***Iterative algorithm to insert list element at tail***

- Linking a new list element to the end of a list can be achieved by an **iterative algorithm** as follows:

We must make a distinction between inserting into an **empty** list and a **non-empty** list....

- If the list at head is **empty**:

```

elem.next = null;
head = elem; (new head)

```

- If the list at head is **not empty**:

```

ptr = "Find the last element in the list at head";
ptr.next = elem;
elem.next = null; (head remains unchanged)

```

- The following is a Java program containing an **Insert** function that inserts the "newelem" list element at

the tail of a linked list that begins at "head":

```

static ListElement Insert(ListElement head, ListElement newelem)
{
    ListElement ptr; // local variable to run down the list

    if (head == null)
    { // Empty list, make "newelem" the first element
        newelem.next = null; // Mark last element
        return(newelem);
    }
    else
    { // find the last element
        ptr = head;
        while (ptr.next != null)
            ptr = ptr.next;
        ptr.next = newelem; // Link new to the last element
        newelem.next = null; // Mark last element
        return(head);
    }
}

```

This function will return the **head** of the new list (with the "newelem" inserted).

---

- Since we are **discussing recursion**, I will **not** spend time discussing the ***iterative algorithm***

But I will **show you** the **code** -- you can study it on your own.

The following program fragment shows how to use this function:

```

ptr = new ListElement(); // Create a new list element
.... (initialize the list element)

head = Insert(head, ptr); // Insert new element into list

```

- Program in Java with **iterative algorithm** to link at tail of list: [click here](#)
  - Program in M68000 assembler with **iterative algorithm** to link at tail of list: [click here](#)
- 
- 

- **A note on allocating memory space for a new list element**

- In **Java**:

```

ListElement ptr;

ptr = new ListElement(); // Create a new list element

```

Recall that the **new** operator will:

- **Reserve** (= allocate) some (enough) **memory space** to store a **ListElement** typed **object**
- And **return** the **location** (= address) of the **reserved memory space**

- 
- To **simulate** the **effect** of the **new** operator, I have **written** the **malloc()** **function** that does the **following**:

- The function **malloc** takes one **input parameter** in **D0**
  - **D0** = the **number of bytes** of **memory** you want to **reserve**
- The **malloc** (memory allocate) function will **reserve** the amount of **bytes** of memory given in **D0**.
- The function **malloc** then returns the **location** (**address**) of the start of reserved memory in register **A0**.

- 
- **Example:**

```
Java:  
  
ListElement ptr;  
  
ptr = new ListElement();  
ptr.value = 1234;  
  
  
M68000 equivalent:  
  
// ptr = new ListElement();  
move.l #8, d0  
jsr    malloc  
move.l a0, ptr  
                reserve 8 byte for a List object  
                put address in ptr  
  
// ptr.value = 1234;  
movea.l ptr, a0  
move.l #1234, (a0)
```