
Second Recursive Function: Fibonacci

- The *classic* Fibonacci function

- The **classic** Fibonacci function

```
int fib(int n)
{
    if (n = 0)
        return 1;
    else if (n == 1)
        return 1;
    else
    {
        return fib(n-1) + fib(n-2);
    }
}
```

- Calling the Fibonacci function

- The **Fibonacci** function is **called** with a statement that look like this:

```
int n, result;
result = fib(n);
```

- Passing parameter **n** from **main** program to **fib()**:

- Because the **Fibonacci** function is *recursive*, each invocation must has its **own** copy of parameter variables
- This can *only* be realized by using a **stack** !!!

Therefore, the **main** program **must** pass the parameter **n** to **Fibonacci** by pushing **n** onto the system stack:

```
move.l n, -(a7)
```

This **instruction** will create the following **stack structure**:

```

+-----+ <----- Stack pointer (A7)
|   parameter n   |
+-----+
|   .....   |
|  rest of the stack  |
|   .....   |

```

- Next, the **main** program will **call** the **fib() function** with a **bsr** instruction:

```
bsr fib
```

This will **push** the **return address** on the **stack** and create the **following stack structure**:

```

+-----+ <----- Stack pointer (A7)
|   return address   |
+-----+
|   parameter n   |
+-----+
|   .....   |
|  rest of the stack  |
|   .....   |

```

- **The Prelude of the Fibonacci function**

- If you **look** at the **Fibonacci function** *carefully*:

```

int fib(int n)
{
    if (n = 0)
        return 1;
    else if (n == 1)
        return 1;
    else
    {
        return fib(n-1) + fib(n-2);
    }
}

```

You will **notice** that there are **two** (**recursive**) calls to **Fibonacci**.

- **Fact:**

- **In order** to *compute*:

$\text{fib}(n-1) + \text{fib}(n-2)$

we **must**:

- **Call** the function $\text{fib}(n-1)$
(and **obtain** the **return value** (x))

- **Then, call** the function $\text{fib}(n-2)$
and **obtain** the **return value** (y)

And add x and y

- **Very important fact:**

- **After** we **obtained** the **return value** (x) from $\text{fib}(n-1)$, we **cannot** compute the **result yet** !!!

- We must **still** find the **second Fibonacci value** ($\text{fib}(n-2)$) **before** we can compute the **result** !!!

Conclusion:

- We must **save** the **return value** (x) in a **safe place** !!!

Where is a **safe place** in **recursive programming**:

- The **only safe place** in to **recursive programming** to **store** values is:
 - **Local variables** on the **stack** !!!

Therefore:

- We **must** create **one local variable** to **save** the **return value** (x) of the call $\text{fib}(n-1)$

- The **prelude** of the **Fibonacci function** is **therefore**:

```
***** PRELUDE
```

```

    move.l a6, -(a7)    ; Save caller's frame pointer
    move.l a7, a6      ; Setup my own frame pointer
    suba.l #4, a7      ; Allocate space for local variable for fib(n-1)
*****

```

I will explain what each one does below. Make sure that you realise that the structure of the stack frame is like this when the prelude is **always** executed:

```

+-----+ <----- Stack pointer (A7)
| return address |
+-----+
| parameter n    |
+-----+
| .....         |
| rest of the stack |
| .....         |

```

- `move.l a6, -(a7)`

This will save the frame pointer on the stack, creating this partial stack frame structure:

```

+-----+ <----- Stack pointer (A7)
| saved a6    |
+-----+
| return address |
+-----+
| parameter n  |
+-----+
| .....       |
| rest of the stack |
| .....       |

```

- `move.l a7, a6`

This will make the frame pointer A6 points to the stack frame that is now being built:

```

+-----+ <---- Frame pointer A6 & Stack pointer (A7)
| saved a6    | point to the same location....
+-----+
| return address |
+-----+
| parameter n  |
+-----+
| .....       |
| rest of the stack |
| .....       |

```

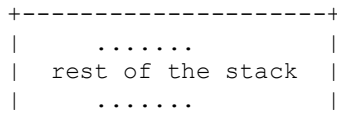
- `suba.l #4, a7`

This will push the stack pointer A7 **4 bytes** up, allocating 1 integer variables -- used to save the **return value** of `fib(n-1)`.

```

+-----+ <---- Stack pointer (A7)
| help (local var) |
+-----+ <---- Frame pointer (A6)
| saved a6    |
+-----+
| return address |
+-----+
| parameter n  |

```



- When the prelude is finish, the stack frame is complete and the actual function can begin.

- **How to access the parameter and local variables inside the `fib` function**

- **How to access the **parameter** and the **local variables** inside the `fib()` function:**

- Parameter **n** is located 8 bytes **below** starting from the address contained in the frame pointer A6.

So the address mode that will let you get to this variable is **8(A6)**

- Local variable **help** is located 4 bytes **above** starting from the address contained in the frame pointer A6.

So the address mode that will let you get to this variable is **-4(A6)**

- **Calling `fib()` from within `fib()`:**

- **How **Fibonacci** calls itself:**
- **Fact:**

- It is the **same way** as **how** the **main program** calls the **Fibonacci function** !!!

- **Method:**

- **Pass** the **parameter** on the **stack**

- **Call** the **Fibonacci** function

- **Clean up** the **parameter**

- **Use** the **return value** (in **D0**)

Make sure you pop the parameter from the stack after Fibonacci returns - because the parameter has not been cleaned up.

The following is the program fragment where Fibonacci calls fib(n-1):

```

    move.l 8(a6), d0    ; retrieve parameter n into register d0
    sub.l #1, d0       ; d0 = n - 1
*
* ----- ; fib is calling fib now !!!!
*
    move.l d0, -(a7)   ; Push (n-1) as parameter
    bsr   fib         ; Call fib(n-1)
    adda.l #4,a7      ; Clean up parameter from stack

    move.l d0, -4(a6)  ; help = return value of fib(n-1) in register D0

```

Fibonacci will call itself a second time with value n-2. The following is the program fragment where Fibonacci calls fib(n-2):

```

    move.l 8(a6), d0    ; retrieve parameter n into register d0
    sub.l #2, d0       ; d0 = n - 2
*
* ----- ; fib is calling fib again...
*
    move.l d0, -(a7)   ; Push (n-2) as parameter
    bsr   fib         ; Call fib(n-2)
    adda.l #4,a7      ; Clean up parameter from stack

    add.l -4(a6),d0    ; Compute the value: fib(n-1)+fib(n-2)

```

I have highlighted the difference between the first call and the second. The second call uses a different parameter value and stores the return value in a different local variable !

- The **full** assembler program:

```

* =====
* main: result = fib(n)
* =====

Start:
    movea.l #12345,a6    ; Store something in a6 to make it dramatic

    move.l n, -(a7)     ; Call fib(n)
    bsr   fib
    adda.l #4,a7        ; pop parameter off the stack
    move.l d0,result    ; result = return value

Stop:  nop

n:     dc.l 5           ; variable n (input)
result: ds.l 1         ; variable result (output)

* ===== Fib
* int fib(int n)

```

```

* {
*   if (n = 0)
*     return 1;
*   else if (n == 1)
*     return 1;
*   else
*     {
*       return fib(n-1) + fib(n-2);
*     }
* }
*
* -----
* Input: n on stack
* Output: fib(n) in register d0

fib:
***** PRELUDE
    move.l a6,-(a7)      ; Save caller's frame pointer
    move.l a7,a6        ; Setup my own frame pointer
    suba.l #4,a7        ; Allocate space for local var. "help"
*****
* Start of function....

    move.l 8(a6),d0     ; n
    cmp.l #0,d0        ; n == 0 ?
    bne  else1

    move.l #1,d0       ; return(1)

***** POSTLUDE
    move.l a6,a7        ; Deallocate local variable(s)
    move.l (a7)+,a6     ; restore caller's frame pointer
*****
    rts

else1:  move.l 8(a6),d0     ; n
        cmp.l #1,d0        ; n == 1 ?
        bne  else2

        move.l #1,d0       ; return(1)

***** POSTLUDE
    move.l a6,a7        ; Deallocate local variable(s)
    move.l (a7)+,a6     ; restore caller's frame pointer
*****
    rts

else2:
***** fib(n-1)
    move.l 8(a6),d0     ; n
    sub.l #1,d0        ; n - 1
    move.l d0,-(a7)    ; Push (n-1)
    bsr  fib            ; call fib(n-1) - will return to next instruction

    adda.l #4,a7        ; Clean up: Pop parameter (n-1) from stack

    move.l d0,-4(a6)   ; Save return value (fib(n-1)) in local var !!

***** compute fib(n-2)
    move.l 8(a6),d0     ; n
    sub.l #2,d0        ; n - 2
    move.l d0,-(a7)    ; Push (n-2)
    bsr  fib            ; call fib(n-2) - will return to next instruction

    adda.l #4,a7        ; Clean up: Pop parameter (n-2) from stack

```

```
*****
        add.l -4(a6),d0      ; Compute the return value: fib(n-1)+fib(n-2)
*****
***** POSTLUDE
        move.l a6,a7        ; Deallocate local variable(s)
        move.l (a7)+,a6     ; restore caller's frame pointer
*****
        rts
End:
        end
```

- **Example Program:** (Demo above code)

Example

- Prog file: [click here](#)

How to run the program:

- **Right click** on link and **save** in a scratch directory
- To compile: **as255 fib**
- To run: use **Egtapi m68000**