Implementing a Stack

- A stack is a data structure with two operation:
  - void Push(int x): put the value "x" in the stack. The value "x" is put "at the top" of the stack.
  - $\circ$  int Pop(): remove the top value off the stack. The value removed is returned.

Example:

<b></b>		
(1) Initially:	++	(empty stack)
(2) After Push(4):	++   4   ++	
(3) After Push(9):		(new value get put ON TOP)
(4) After Push(1):		(new value get put ON TOP again)
Suppose we have 3 vari	ables "a", "b" a	and "c" defined
<pre>(5) After a = Pop() and a = 1</pre>		(top value is REMOVED !)
(6) After b = Pop() and b = 9	: ++   4   ++	(top value is REMOVED again !)
(7) After $c = Pop()$ and $c = 4$	: ++	(Stack is empty)

- Stack in a program versus the System Stack
  - The stack that you learned in Java is usually implemented using an array variable

Example: (CS171 material)

```
public class Stack
{
   int[] A;
                      // Points to the top of the stack
   int stacktop;
  public Stack(int size)
                         // Create the array to hold values in stack
      A = new int[size];
      stacktop = -1;
  public void push( int x )
   ł
                              // Move stack top and put x on the stack
      A[++stacktop] = x;
   }
  public int pop( )
   ł
      return A[stacktop--];
   }
3
```

• When a **program** is **running**, the **computer system** will need to **store** various **information items** (variables) in a *stack* 

**Therefore:** 

All modern computers maintain a program stack (a.k.a. System stack) to manage information

Information stored in the system stack include:

• function activation information (created when a function is called

(These function activation information will be removed (popped) when a function returns (exits))

- The reason we want to use a stack is the order the elements enter and leave the stack is exactly the same order of function activation and deactivation.
- The system stack is (always) stored inside the memory (memory stores everything :-))
- When a program is running, the memory is organized as follows:

```
Memory:
   0 +----
                  <--- this part contains the "code"
          Α
                        This part of the memory is "write protected"
                        (can not - and should not - be modified)
                 <--- this part contains "static variables"
                       These variables exist at the start of the
         в
                1
                       program (a.k.a.: compiled-time variables)
                |\ | <--- this part grows when the program creates
                         new objects: reserve space for instance
                1 1
         С
                         variables in the new object
                1 1
                l v
                v <--- Direction of growth is "downwards"</pre>
                1 77
                | <---- this part is "free" memory</pre>
                        (Free means: unreserved !)
                | ^ <--- Direction of growth is "upwards"
                1.1
         D
                1 1
                \mid \mid <--- This part grows when the program invokes
       System | |
                       a function/method, reserve space for:
        stack
              (1) return address
                (2) parameter variables
                | v
                       (1) local variables
     +----+
1. Parts A and B exist as soon as the program starts execution
  and will exist throughout the program execution.
2. Part C will grow when program creates new objects
   (with the new operator in Java or malloc() in C)
  Part C is called the "system heap"
3. Part D will grow when program invokes a function/method
   Part D is called the "system stack"
```

• Implementing the System Stack

• Just like a stack object in Java, in order to implement a stack, we need:



• *Unlike* in a programming language, we have access to the entire computer memory when we program in assembler !!!

So the system stack that is located at the *end* of the memory area:

Memory:	
0 +     A	+   < this part contains the "code"   This part of the memory is "write protected
   +	(can not - and should not - be modified)   +   < this part contains "static variables"
B   +	These variables exist at the start of the program (a.k.a.: compiled-time variables)
l I C	<pre>  ^   &lt; this part grows when the program creates    </pre>
 +==============	▼ < Direction of growth is "downwards"   ▼ +
	< this part is "free" memory         (Free means: unreserved !)   
+=====================================	+   ^   ^ < Direction of growth is "upwards"
	^         < This part grows when the program invokes
System   stack 	<pre>    a function/method, reserve space for:     (1) return address     (2) parameter variables   v (1) local variables</pre>

do not need to be *defined* (as an array as in Java)

(The **memory** is **there** and it is *reserved* for us to make a **stack**)

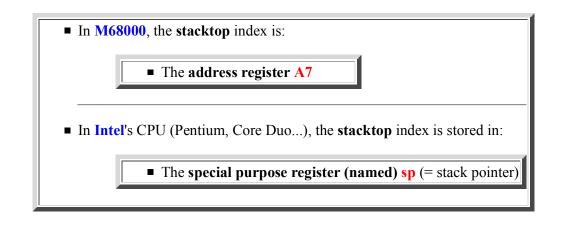
• All we need to **implement** the **System Stack** is:

A stacktop index

• The system stack is always implemented by:

• Using *one* of the registers in the CPU as stacktop index

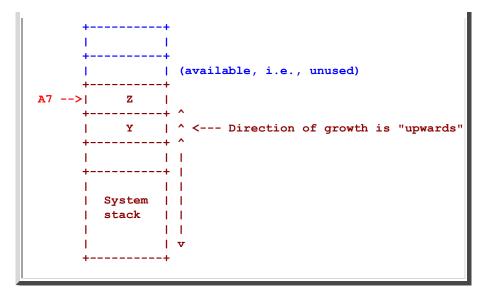
**Example:** 



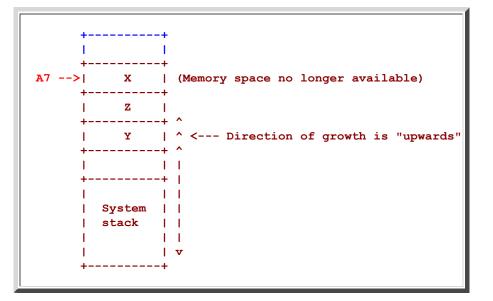
• So in M68000, the system stack is given as follows:

Memory:	
0 +	+   < this part contains the "code"   This part of the memory is "write protected"   (can not - and should not - be modified) 
   B 	+   < this part contains "static variables"   These variables exist at the start of the   program (a.k.a.: compiled-time variables) +
c	<pre>^</pre>
+=========	+
	< this part is "free" memory        (Free means: unreserved !)   
A7>+===================================	+   ^   ^ < Direction of growth is "upwards"   ^
D	
System	<pre>    &lt; This part grows when the program invokes     a function/method, reserve space for:</pre>
System	a function/method, reserve space for:     (1) return address
	(2) parameter variables
	v (1) local variables
+	+

- Pushing values onto the System Stack
  - Suppose the System Stack is as follows before we push a new value:



• The **Suppose** the **System Stack** will be **like this** *after* we **push** a new value *X*:



(The **memory space** is **no longer available** because **all memory spaces** "at and below" the **stack pointer A7** are *used* !!!!)

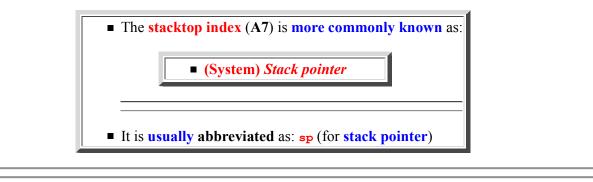
• The following assembler instructions will push a integer (4 byte) value X onto the System Stack:

suba.1#4, A7// Move system stacktop up 4 bytesmove.1X, (A7)// Store X in the top of the stack

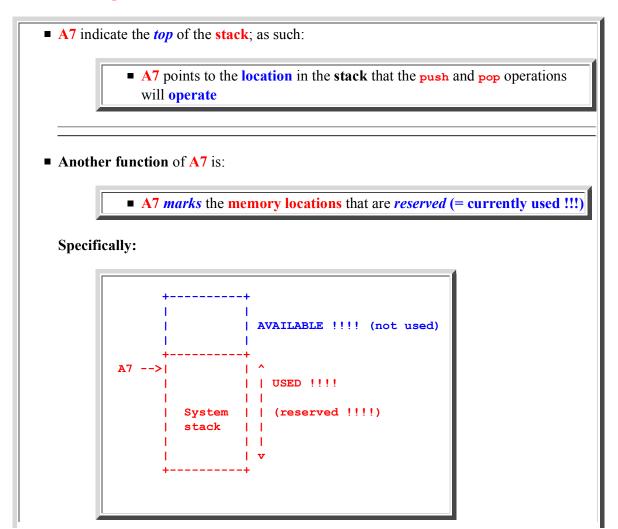
- De-allocating variables from the System Stack
  - The following assembler instructions will destroy (= unreserve memory)) a integer (4 byte) value from the System Stack:

a	dda.1	#4, A7	// Move	system stacktop down 4 byte	≥s

- The system stack pointer
  - Stack pointer:



- Important functions of the system stack pointer A7
  - Function of the stack pointer A7:



• Therefore:

